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### Type theory and Universal Grammar

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### Universal Grammar

Introduction

The idea of Universal Grammar (UG) as the hypothetical linguistic structure shared by all human languages harkens back at least to Roger Bacon in the 13th century [Ranta, 2006]

- The modern notions of UG:
  - Substantive UG [Chomsky, 1970, Chomsky, 1981, Chomsky, 1995]
  - "Diluted" UG: the Language Acquisition Device [Jackendoff, 2002]

### Universal Grammar

Introduction

- The ideas of UG occur in the broader context of...
  - Substantive universals [Plank et al., 2009]
  - Implicational universals [Greenberg, 1966]

Greenberg, J. H. (1966). Some universals of grammar with particular reference to the order of meaningful elements. In Greenberg, J. H., editor, Universals of Grammar, pages 73–113. MIT Press, Cambridge, MA.

Plank, F., Mayer, T., Mayorava, T., and Filimonova, E. (2009). The Universals Archive. https://typo.uni-konstanz.de/archive

Introduction

(0) type := a category of semantic value.

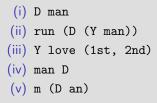
By (0), type theory is by definition suited for analyzing universal phenomena in natural language (NL), as NL semantics in largely universal (as witnessed by the possibility of translation from any human language to another)

Introduction

Thus, if we could build a fundamentally semantic description of grammar (e.g. one on top of and integrated with a semantically universal description of NL), it might at least stand a chance of being universal

Interpreting NL in

#### NL expressions as function applications:



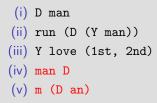
Complex formulas are written in prefix notation, a b or a(b), with a standing for a function and b for argument(s)

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Left-associativity, i.e. left to right valuation

Interpreting NL in

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Interpreting NL in

- Call the formulas (i)-(iii) formulas of a specification language (SL)
- Then, we specify SL formulas from NL expressions and derive NL expressions from SL formulas

Axioms

- (1) Arguments must be either specified or derived before the relation expressions in which they appear
- (2) NL and SL expressions are well-formed both syntactically and semantically, i.e. well-formed and well-typed
- (3) For a particular language, the symbols are type constants; in UG they are type variables (e.g. man valuates to *man* in English and *homme* in French)

in  $\mathsf{Coq}^\dagger$ 

(vi) Y know i (who (Y COP ill (the man)))
(vii) Y know i (who (Y COP (ill, the man)))
(viii) Y know i (who (Y COP (ill, D man))) : S : U
(ix) PRES know i (who (PAST COP ill (the man)))

where S is sentence,  $\mathcal{U}$  the top-level universe in SL, and "x : y" := "x has type y"

<sup>†</sup>https://gitlab.com/eluuk/nlc/blob/master/cop.v

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Optional

- Selectional restrictions are (onto)logical restrictions on the types of arguments of NL relations (e.g. [Asher, 2014, Luo, 2010])
- For example, an adjective like *red* imposes the restriction that its argument be of type physical entity, while a verb like *know* imposes a restriction that its subject be a sentient and object an informational entity

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#### Optional

$$\begin{array}{ll} (\mathsf{x}) \; [\texttt{red}] : \; X_{\mathsf{Phy}} \to \mathsf{P} \\ (\mathsf{xi}) \; [\texttt{know}] : \; X_{\mathsf{Sen}} \to Y_{\mathsf{Inf}} \to \mathsf{S} \end{array}$$

where [x] is the interpretation of x, P is phrase, S sentence and X, Y are type variables indexed by selectional restrictions

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The rule of metaphor or metonymy elimination

$$\frac{u_j: Z_j \qquad (x_h^e: X_h^e) \mapsto (y_j^e: Y_j^e)}{x \dots (u_j)^e \dots : W}$$
 MM-Elim,

where  $x_h^e$  is a function x, eth argument of which is restricted to h;  $X_h^e$  a function type X, eth argument type of which is indexed by h;  $u_j$  a (possibly empty) argument u, restricted to j;  $x \mapsto y := "x$  is interpreted as y"; and  $X_h^e$ ,  $Y_j^e$ ,  $Z_j$ , W : U, where Uthe top-level universe in SL

The rule of metaphor or metonymy elimination

- By MM-Elim, whenever we have a metaphor/metonymy (x<sup>e</sup><sub>h</sub> is interpreted as y<sup>e</sup><sub>j</sub>) and possibly u<sub>j</sub>, x ... (u<sub>j</sub>)<sup>e</sup>... is well typed in SL (and NL)
- For example, {idea<sub>Inf</sub>, (red<sub>Phy<sup>1</sup></sub> → communist<sub>Inf<sup>1</sup></sub>)} ⊢ red idea<sub>Inf</sub> : W
- As we take all elementary arguments to be nullary relations, we also have {red<sub>Phy<sup>0</sup></sub> → communist<sub>Inf<sup>0</sup></sub>} ⊢ red<sub>Inf</sub> : W for argumental uses of the words

The problem

- The received view among typological linguists is along the lines that nothing in NL is universal [Haspelmath, 2007, Evans and Levinson, 2009]
- But what about sign, form, meaning, word, sentence, morpheme, phrase, etc.?
- Besides these general counterexamples, the main difficulty is conceptual rather than factual, being due to the virtual non-existence of universally shared definitions

A proposal: Universal categories

We propose some universal linguistic categories, defined by their function:

proper name (PN)
connective (CON — and, but, or, not...)
XP (Frequently also referred to as NP or DP)
declarative sentence (S — john is here...)
interrogative sentence (IS)
connective composition (CONC — x and y, x or y, not x...)

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 But this is about as far as conventional grammatical categories get us

Another proposal: Universal supercategories

#### We propose some universal linguistic supercategories, defined by their function:

- case/adposition (CA nominative, accusative, to, from...)
- case/adposition phrase (CAP john, him, to the house)
- numerals/quantifiers (Q all, some, no, few, one, two...)
- determiner/demonstrative (D a, the, this, those...)
- tense/aspect/mood/voice (TAM the canonical verbal inflection)
- adverbs or other adverbial phrases (ADL quickly, in a hurry...)

#### **Flexibles**

And such

- To proceed with defining universal supercategories, we use the general polymorphic linguistic category of flexibles [Luuk, 2010], exemplified by words like *sleep* and *run* in English
- Since sleep and run "flex" between relation and argument, they are flexibles-over-relation-and-argument.
   Provisionally, we type them X/R (with R for relation and X (from XP) for argument)

 There are many categories of flexibles, e.g. have is a flexible between auxiliary verb (AUX) and infinitival relation (IR) (has type AUX/IR)

Core relation and argument

- Now we can posit the universality of the following supercategories:
  - core relation (R verb, copula, infinitival relation, auxiliary verb or flexible-over-relation)
  - core argument (X noun, proper name, pronoun, gerund or flexible-over-argument)
- Examples: an infinitival relation is *like* in *i like to run*, an auxiliary verb is *must* in *i must run* and a gerund is *running* in *running is healthy*

### The Coq tests

With a flexible

Here are some tests of a Coq implementation of the flexible that is polymorphic between function and argument:

```
Check sleep: gs _ _ _ _ . (*typed as argument*)

Check sleep: NF \rightarrow _ \rightarrow S. (*typed as function*)

Fail Check PAST sleep man. (*fails since "man" is not an XP*)

Check PAST sleep (a man). (*a man slept*)

Check PAST sleep (few (PL man)). (*few men slept*)

Check AST sleep (a (few (PL man))). (*a few men slept*)

Check a sleep.

Fail Check PAST sleep (a hut). (*a hut slept: wrong SR*)

Fail Check PAST sleep (a sleep). (*a sleep slept: wrong SR*)
```

### The Coq implementation

#### Of Ds

#### Let us define some Ds:

a:  $\forall \{x \ y \ z \ w\}$ , gs x y z SG w  $\rightarrow$  gp cp\_x y z SG w the:  $\forall \{x \ y \ z \ u \ w\}$ , gs x y z u w  $\rightarrow$  gp cp\_x y z u w this:  $\forall \{x \ d \ w\}$ , gs cs\_s x d SG w  $\rightarrow$  gp cp\_x x d SG w these:  $\forall \{x \ d \ b \ w\}$ , gs cs\_p x d b w  $\rightarrow$  gp cp\_p x d PLR w

The Ds have function types, with arguments in {...} implicit (implicitly applied). gs \_ \_ \_ \_ and gp \_ \_ \_ \_ are some compound types (in this case, function applications), so e.g. a is a function from gs \_ \_ \_ SG \_ to gp cp\_x \_ \_ SG \_. In Coq, \_ marks any admissible term or type, and SG stands for singular, i.e. a takes arguments in singular and returns phrases in singular.

## The Coq implementation

Of universals

In Coq we can also define universal notations, e.g.

```
Parameter D: \forall {x y z u w}, gs x y z u w \rightarrow gp cp_x y z u w.
(*universal "D" declared as a variable*)
Notation D' := (_: gs _ _ _ \rightarrow gp cp_x _ _ _).
(*universal "D" defined as a notation*)
```

The universality of D and D' comes from x, y, z, u, w and \_ standing for any admissible term or type, whence the applicability of D and D' whenever one of a, the, this, these applies (an ex. with the man knows a few men):

```
Check PRES know (the man) (a (few (PL man))).
Check TAM know (D' man) (D (Q (PL man))).
```

### Conclusion

Universal Grammar

- We have shown how to build an extensive and robust substantive UG using supercategories (categories of categories)
- ▶ In particular, we proposed the following approach to UG:
  - ▶ The universality of categories PN, CON, XP, S, IS and CONC
  - The universality of supercategories CA, CAP, Q, D, TAM, POS, ADL, R and X

## Conclusion

#### Type-theoretical modeling

- ► The main contributions:
  - An integrated modeling of syntax, morphology and compositional semantics (in the form of selectional restrictions).
  - To account for systematic violations of selectional restrictions by metaphor and metonymy, we have shown how to model them type-theoretically.
  - An implementation of a fragment of NL and UG in the Coq proof assistant (https://gitlab.com/eluuk/nlc).
  - The implementation shows how to model many (super)categories of NL, some of them universal, in a purely functional type system (i.e. one comprising only functions and their types) with dependent and polymorphic types.
  - It seems likely that the underlying formalism could be also encoded in a simpler type system, which remains a future work.

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